

Seminar on Secure Software

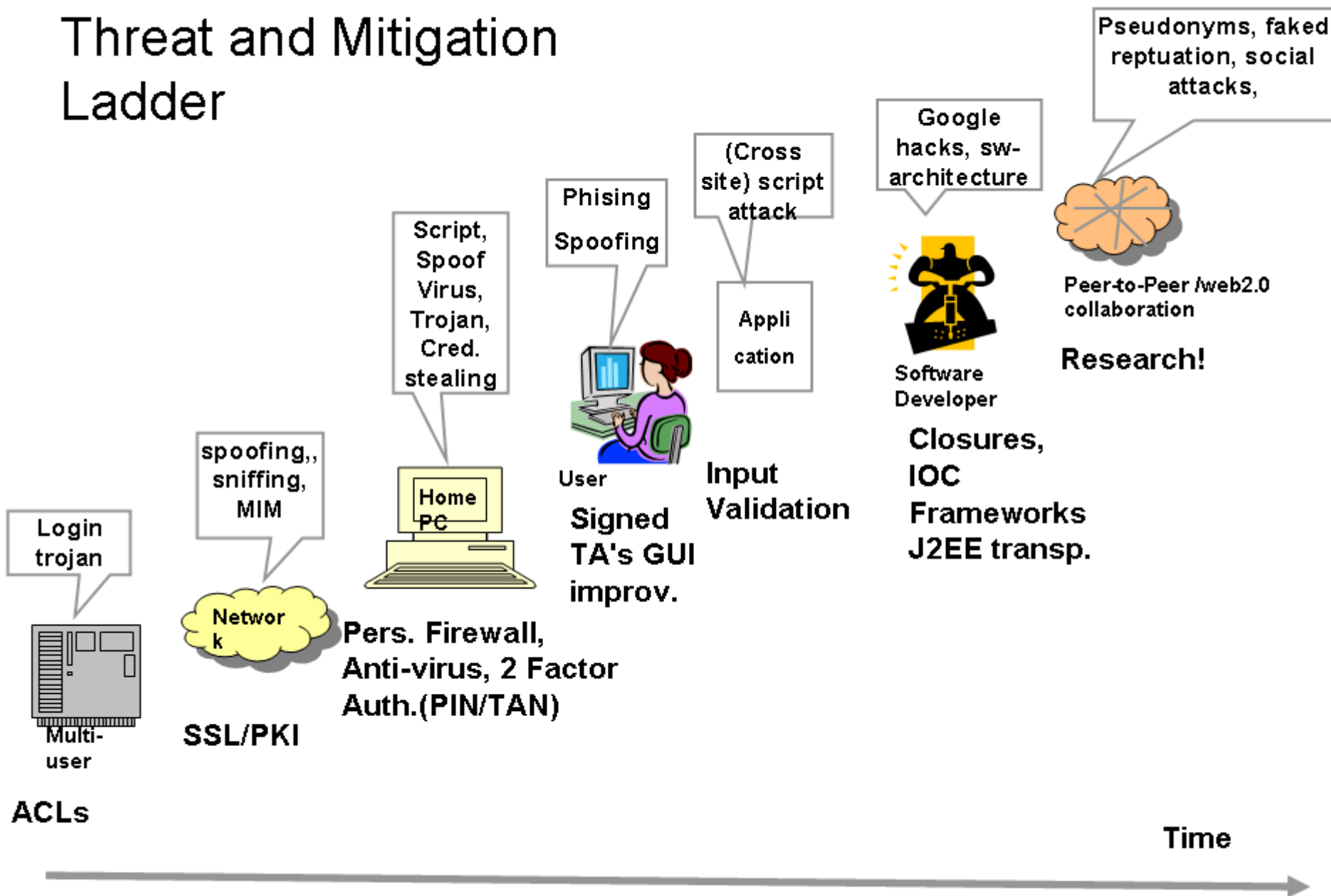
Attacks, Mitigation and fundamental software problems

Input Validation, Filtering and
Damage Control as Software
Mechanisms

Attack Examples

XSS, XSRF, Buffer Overflows,
Character Aliases etc.

Threat and Mitigation Ladder



Input/Output
Related

- A1 [Unvalidated Input](#)
- A4 [Cross Site Scripting](#)
- A5 [Buffer Overflow](#)
- A6 [Injection Flaws](#)
- A7 [Improper Error Handling](#)
- A9 [Application Denial of Service](#)

Infrastructure

- A8 [Insecure Storage](#)
- A9 [Application Denial of Service](#)
- A10 [Insecure Configuration Management](#)

AAA related

- A2 [Broken Access Control](#)
- A3 [Broken Authentication and Session Management](#)
- A9 [Application Denial of Service](#)

System Engineering

- A9 [Application Denial of Service](#)

A "Phishing-Link" to LBBW Bank: XSS due to bad input validation

Hostname of bank:

<http://www.lbbw.de/lbbw/html.nsf/webdokumente/framebooster.htm?OpenDocument&uri=http://www.google.de>

Attack URL (in reality: some IP address or a name close to the original site name like lbbw-systems, lbbw-tech etc.)

Landesbank Baden-Württemberg - Mozilla Firefox

File Edit View Go Bookmarks Tools Help

← → ↻ × 🏠 LBBW <http://www.lbbw.de/lbbw/nml.nsf/webdokumente/framebooster.htm?OpenDocument&url=> Go

Getting Started Latest Headlines Chapter 8. Security ... Technical Tip - JAAS ...

LBBW Landesbank Baden-Württemberg Welcome to the kriha.org homepage

[Personalisierte Startseite](#) | [Anmelder](#)

Google™

Deutschland

Web [Bilder](#) [Groups](#) [Verzeichnis](#) [News](#) [Froogle](#) [Mehr »](#)

[Erweiterte Suche](#)
[Einstellungen](#)
[Sprachtools](#)

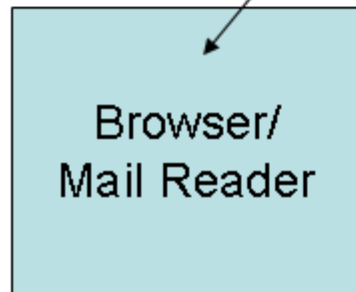
Suche: Das Web Seiten auf Deutsch Seiten aus Deutschland

[Werbung](#) - [Unternehmensangebote](#) - [Über Google](#) - [Google.com in English](#)

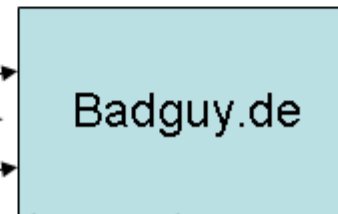
©2006 Google

Phishing Mail: „Dear Customer of mybank...“
 www.mybank.de

1. Trick User into clicking on URL



2. User connects to badguy.de



6. Man-in-the-middle modifies transactions on the fly. Modifies Responses too.

TAN

5. User does Transaktionen

8. User sends TAN to badguy

TAN

3. Badguy forwards requests to bank and sends responses back to user

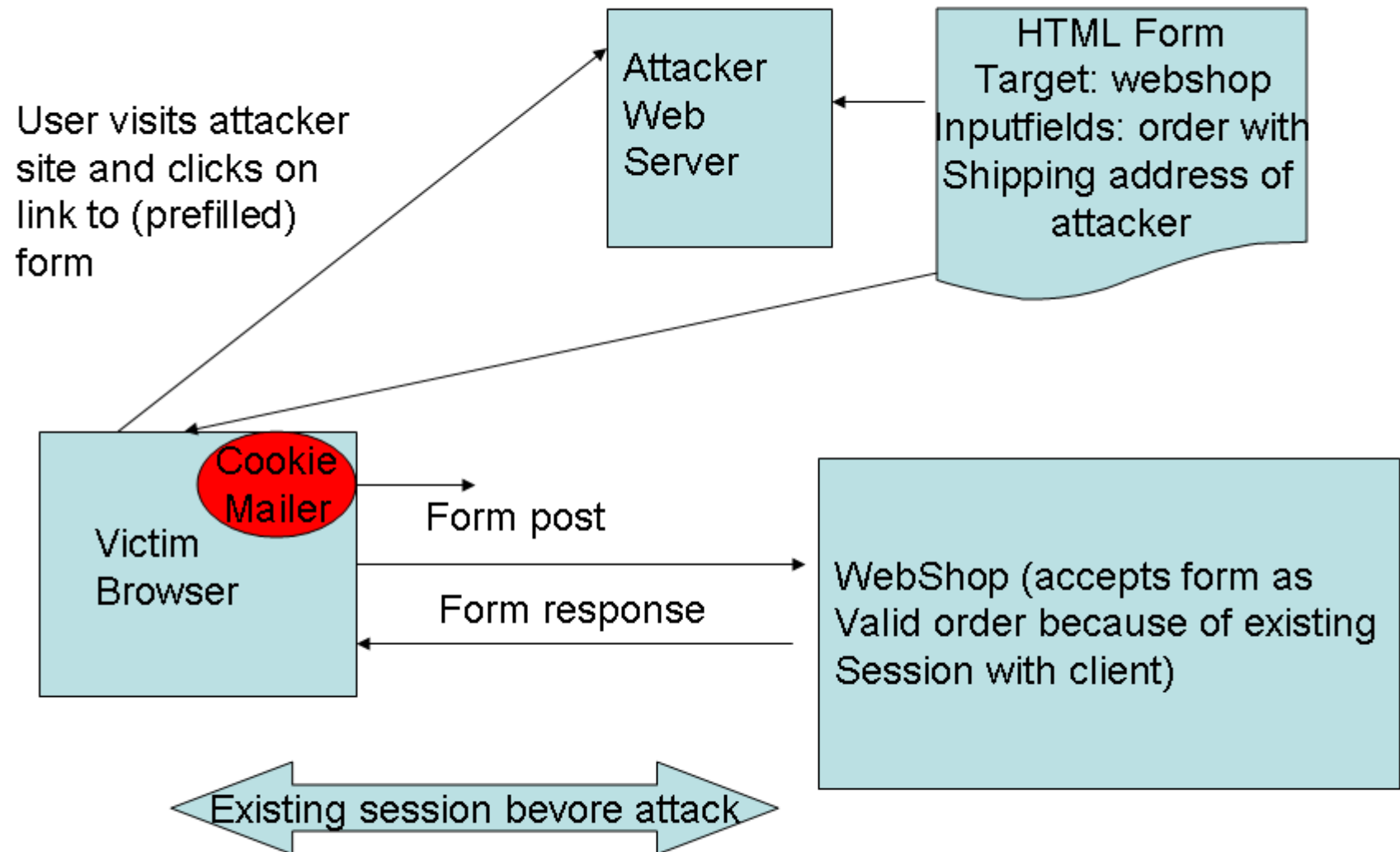
SMS/TAN



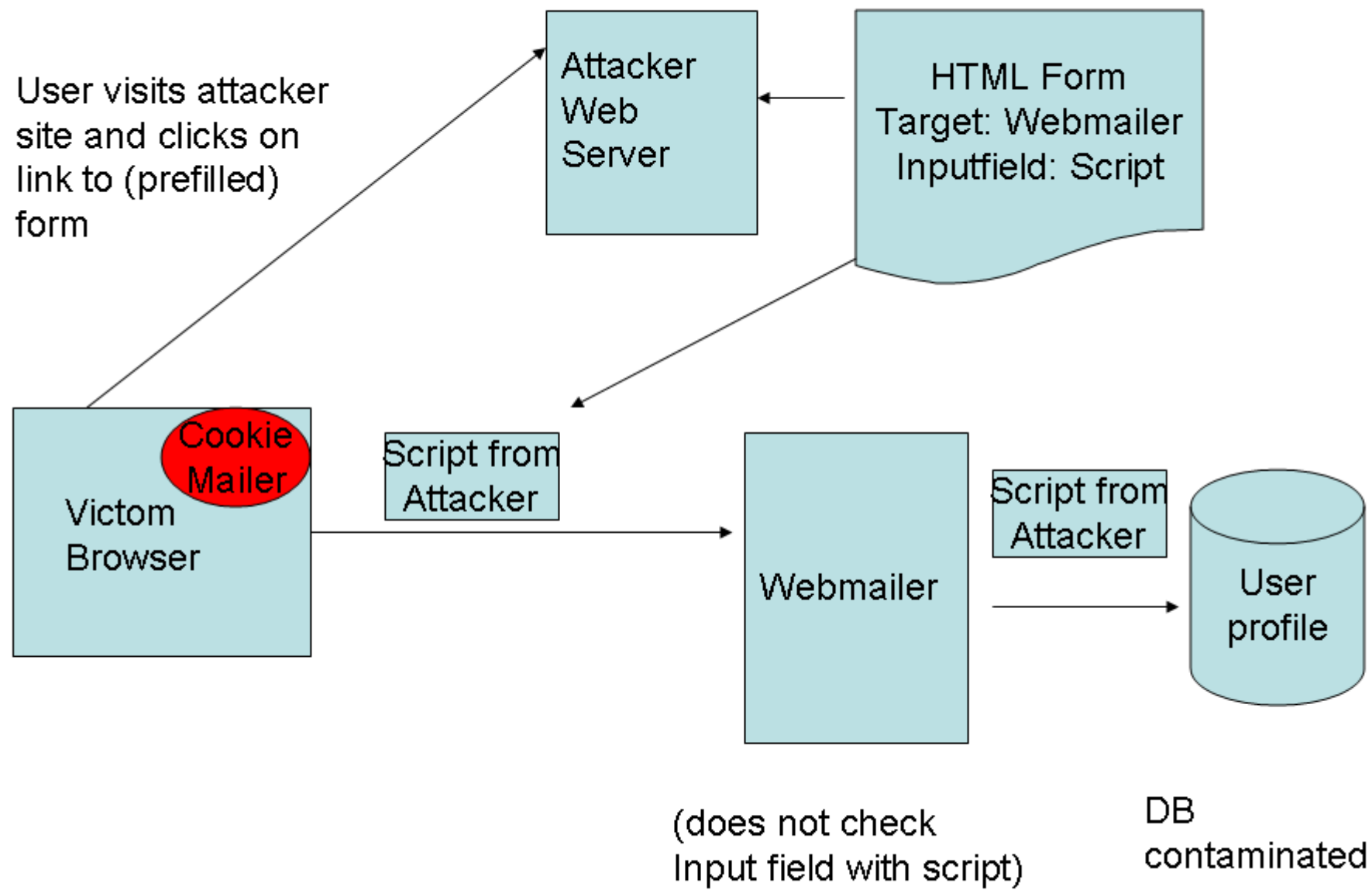
4. Bank asks user to login.

7. Bank sends Users sms with TAN.

Cross-Site Request Forgery (XSRF or Web-trojan)



Injection Attack



```

#include <stdio.h>

int main(int argc, char** argv) {

    int foo=0xeeee;
    char myArray[4];
    gets(myArray);
    printf(" print integer first: %x ", foo);
    printf("%s ", myArray);

}

```

Keyboard Input (with return)	Display Output
a	Eeee a
aa	Eeee aa
aaa	Eeee aaa
aaaa	Ee00 aaaa
aaaaaaaaaaaaaaaa	Core dump with EIP = 6161616161616161 (Hex 61 == `a`)

Our „aaaaaaaa..“ input from keyboard is now the address where the next instruction should be read by the CPU. Now we know how to point the CPU to code we placed on the stack

```
Exception: STATUS_ACCESS_VIOLATION at eip=61616161
eax=00000012 ebx=00000004 ecx=610E3038 edx=00000000 esi=004010AE
edi=610E21A0
ebp=61616161 esp=0022EF08
program=D:\kriha\security\bufferoverflow\over.exe, pid 720, thread main
cs=001B ds=0023 es=0023 fs=003B gs=0000 ss=0023
Stack trace:
Frame  Function  Args
 90087 [main] over 720 handle_exceptions: Exception:
STATUS_ACCESS_VIOLATION
 104452 [main] over 720 handle_exceptions: Error while dumping state
(probably corrupted stack)
```

A program crash is a way into the system!

The kernel trap interface

your code wants to send a message msg to stdout:

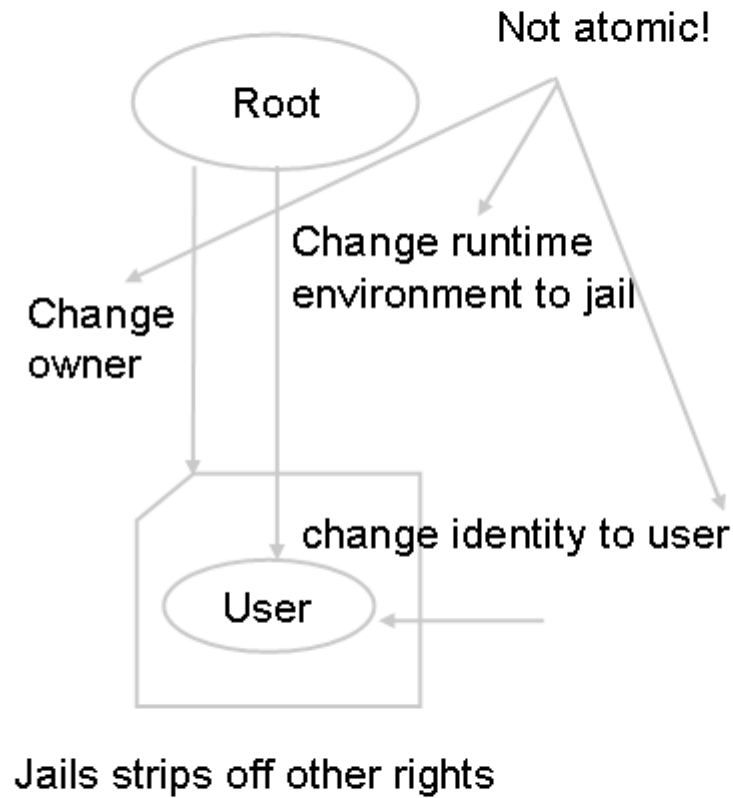
```
push len    ;message length
push msg    ;message to write
push 1      ;file descriptor (stdout)
mov  AX, 0x4    ;system call number (sys_write)
int 0x80      ;kernel interrupt (trap)
add  SP, 12    ;clean stack (3 arguments * 4)
push 0      ;exit code
mov  AX, 0x1    ;system call number (sys_exit)
int 0x80      ;kernel interrupt we do not return from sys_exit there's no need to clean stack
```

The trap (system call interface) is very important for attack code because it is POSITION INDEPENDENT! Your code is NOT LINKED with the running program and therefore does not know where specific library functions etc. are located in your program. The kernel interface is always just there and can be used to load Dynamic Link Libraries into the program.

Attack Vectors on Web Services:

- Wrong input length of variables
- Variables containing wrong characters or meta-characters
- Variables containing SQL commands
- Responses which expose SOAP error codes

Administration and Race Conditions: toc2tou bugs



Admin:

Attacker (knows temp filename)

`Ln -s /etc/passwd /tmp/myFile`

`touch /tmp/myFile`

`... processing...`



Overwrites passwd

`rm /tmp/myFile`



Deletes passwd



Time

SetUid Program:

Victim

Fstat(/tmp/myFile)

Chgrp foo bar



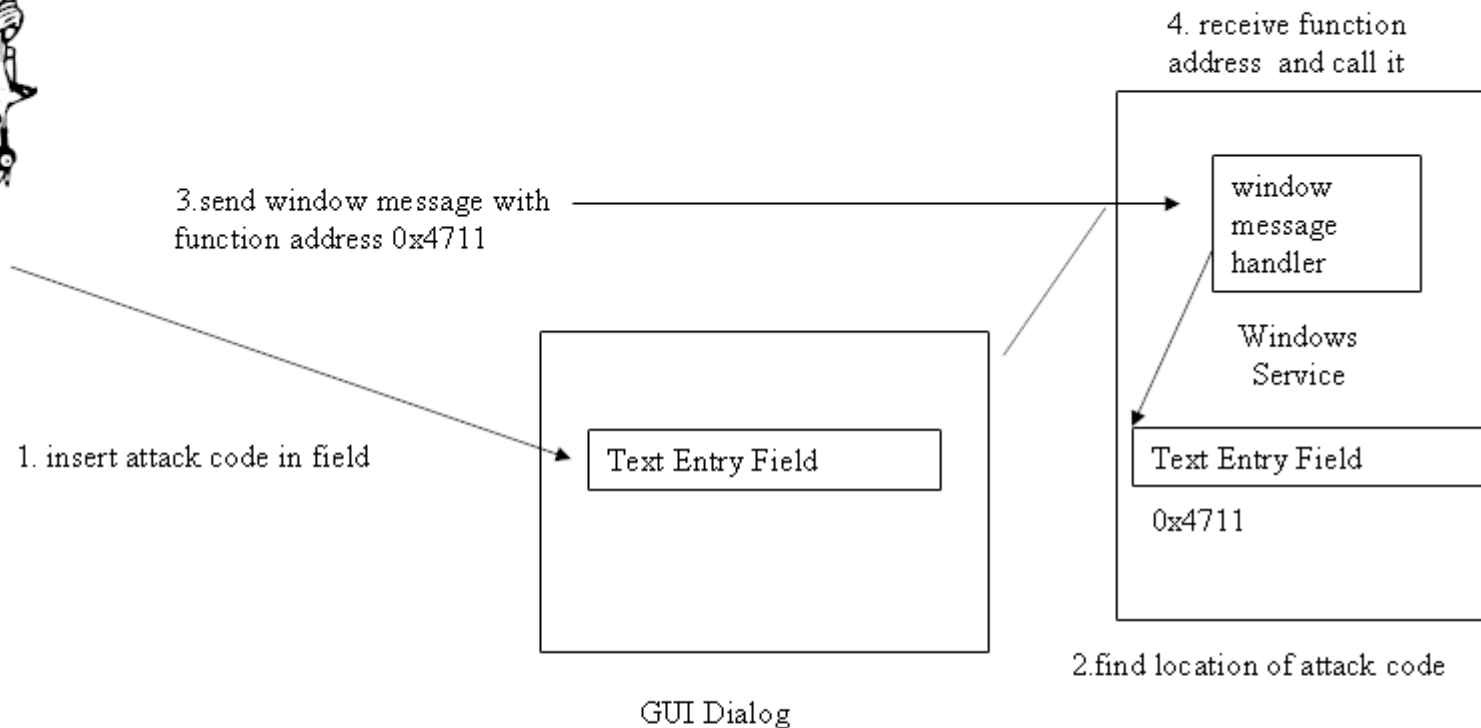
Open(/tmp/myFile)

... processing...



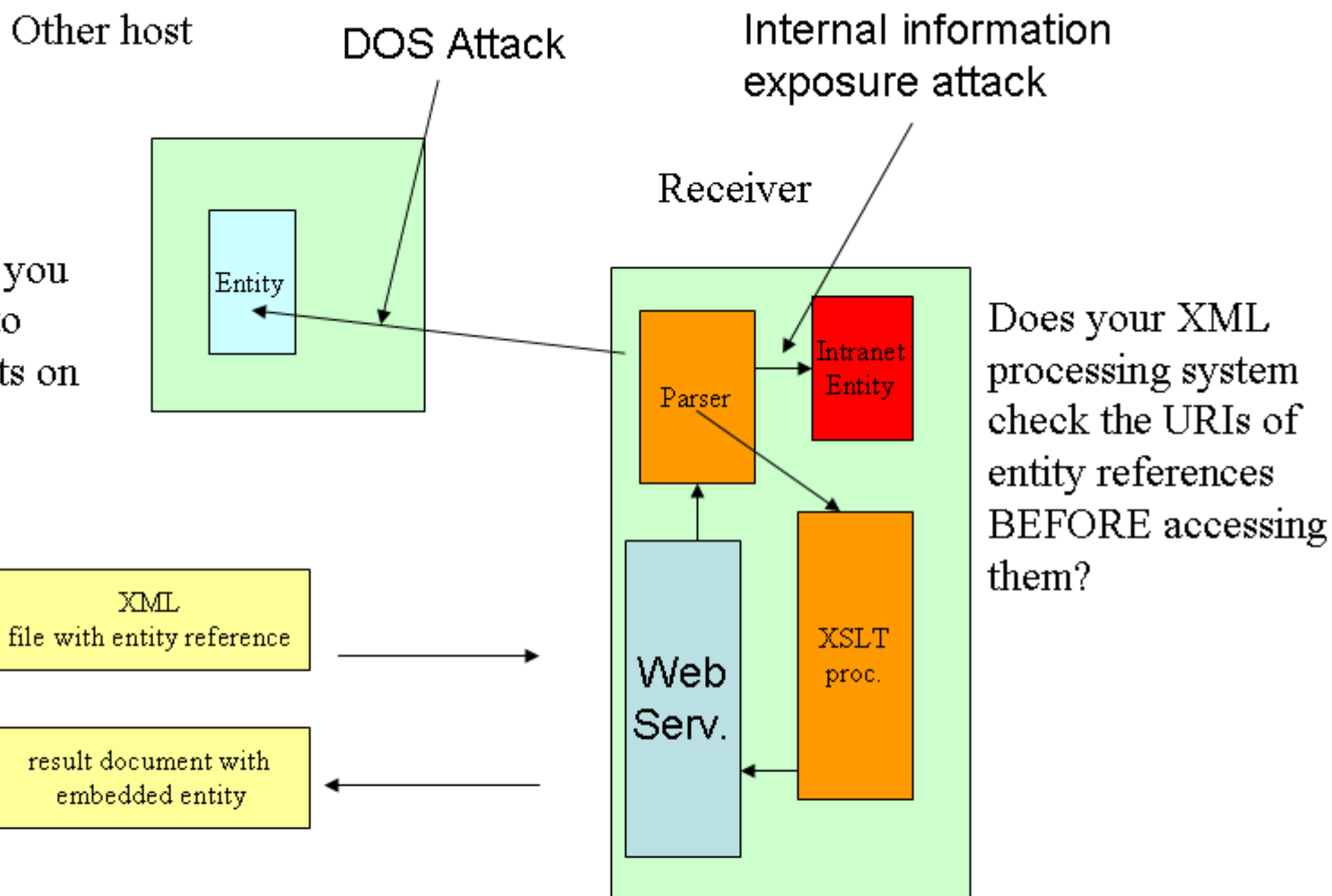
Time

Shatter Attack: fundamental software design flaws



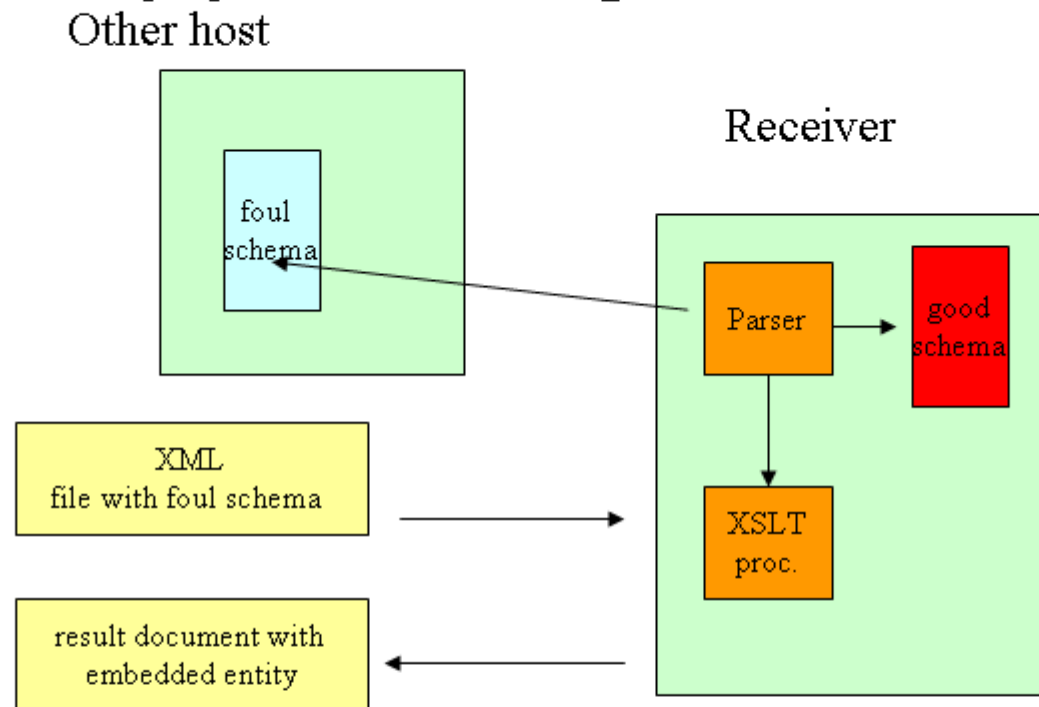
Here the danger is that any program can send certain window messages which contain function addresses IN THE RECEIVERS ADDRESS SPACE. By placing some attack code into the receiver (not hard if a GUI is used by the receiver) the attacker can then direct the receiver message handler to direct control flow to the attack code (step 4 above).

If you offer a rendering service you might be abused to create artificial hits on some host.



```
<?xml version='1.0'?>
<xsl:stylesheet
xmlns:xsl=http://www.w3.org/1999/XSL/Transform version='1.0'>
<xsl:output method="html" encoding="ISO-8859-1" indent="no"/>
<!-- ===== -->
<xsl:script language="java" implements-prefix="sy"
src="java:java.util.system"/>
<xsl:template match="*">
  <xsl:message>
    <xsl:text>No template matches </xsl:text>
    <xsl:value-of select="sy:exec(...)" />
    <xsl:text>.</xsl:text>
  </xsl:message>
```

Suppressing Validation



James Clark mentioned recently an especially evil way to work around validation: „Suppose an application is trying to use validation to protect itself from bad input. It carefully loads the schema cache with the namespaces it knows about, and calls validate(). Now the bad guy comes along and uses a root element from some other namespace and uses xsi:schemaLocation to point to his own schema that that has a declaration for that element and uses `<xs:any namespace="##any,, processContents="skip"/>`. Won't they just have almost completely undermined any protection that was supposed to come from validation?“

Code points for most characters in the languages of the world

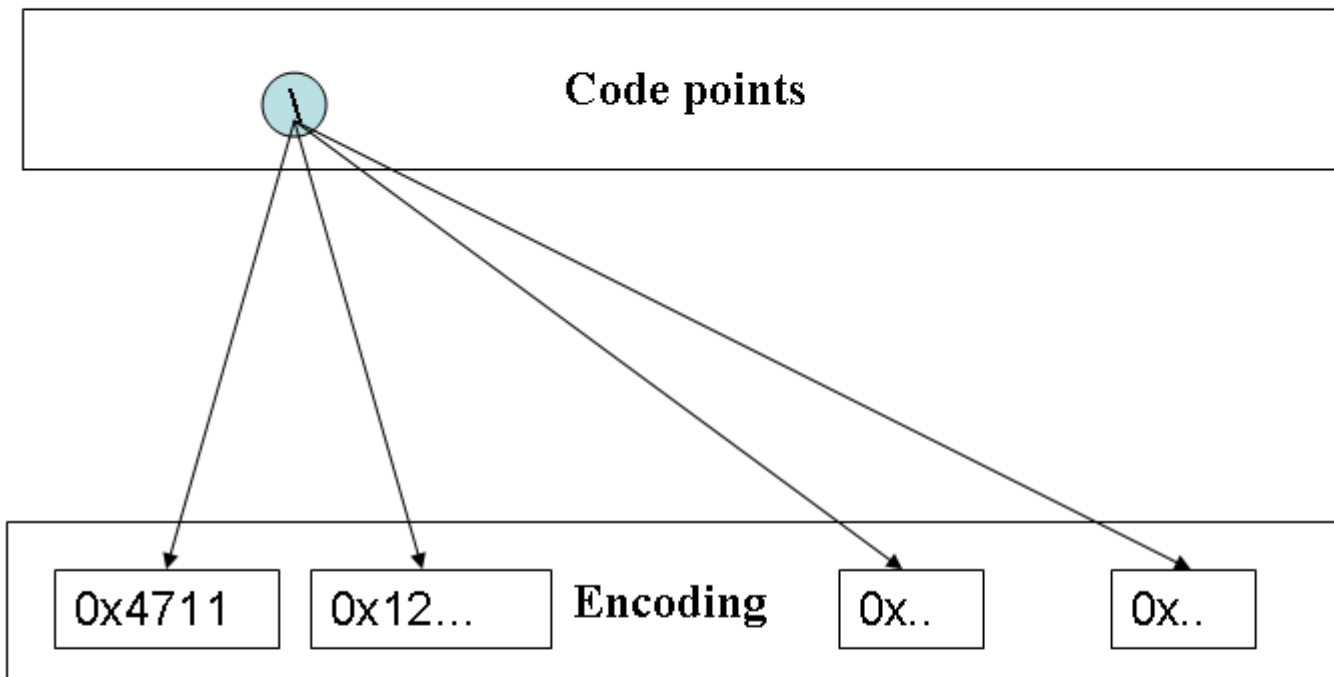
Unicode code points
(names and numbers
of characters) 9% of 4
Gigabyte

**UTF8, UTF16 or UTF32 Encodings of code points
(code units or blocks)**

3 different ways to
encode ALL code
points (size vs.
performance)

arbitrary glyphs (fonts)

Not defined by
unicode.



One codepoint can have several different encodings. Filter code needs to NORMALIZE FIRST and then FILTER!

Filter code to detect ...\ attacks:

```
If (encoded == 0x4711)
```

```
    removeCharacter();
```

```
// what about the other possible encodings of backslash????
```

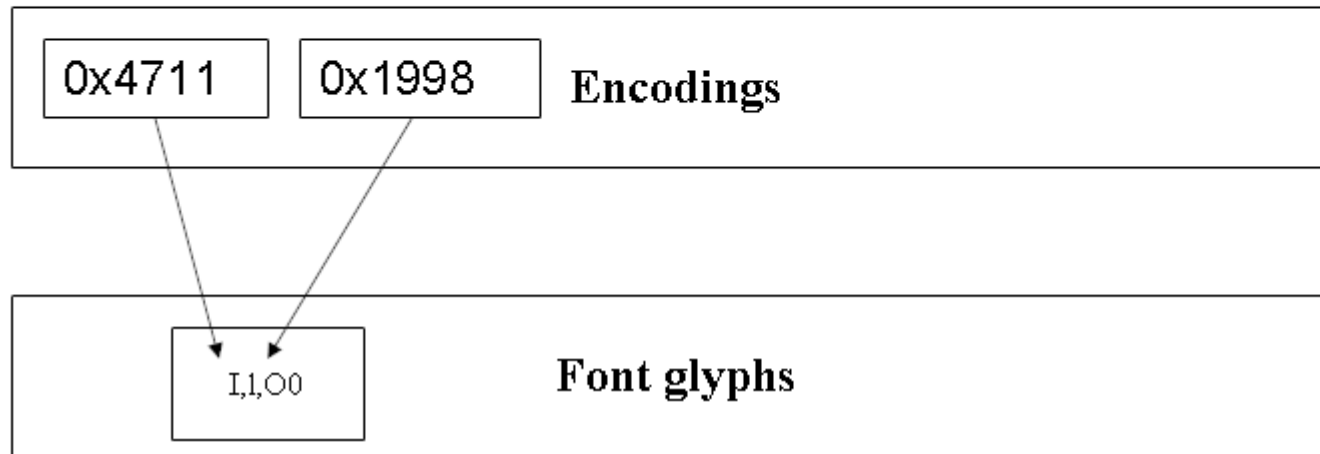
Unicode Exploit

code point U+0000

encoded as: 0, 110 00000 10 000000, etc.

Unicode code points
(names and numbers of
characters) 9% of 4
Gigabyte

Processors are not allowed to interpret any encoding other than the shortest form, in this case 0. Otherwise the extended forms could escape filtering and become active during interpretation.



**Fonts can display
unicode code
points any way
they want.**

One visual „look“ (e.g. lowercase „l“ and uppercase „I“ or greek omicron vs latin o.

Unicode homographs and DNS

Two different code points

ASCII DNS

Unicode Characters DNS

DNS names can now
contain Unicode characters

two different fonts

Not defined by unicode.

l,l,O

One visual „look“ (e.g. lowercase „l“ and uppercase „I“
or greek omicron vs latin o.

The firefox browser switched back to showing the unicode escape sequences in domain names to allow the user to differentiate e.g. a latin „a“ from a kyrillic „a“. Otherwise the user could be tricked into connecting to www.ebay.com with the „a“ being really the kyrillic version. In this case the user would connect to the wrong site. Expect many more security problems with unicode in the future, especially in the GUI area.

[Anmelden](#)



Web [Bilder](#) [Groups](#) [News](#) [Froogle](#) [Mehr »](#)

inurl:jmx-console

Suche

[Erweiterte Suche](#)
[Einstellungen](#)

Suche: Das Web Seiten auf Deutsch Seiten aus Deutschland

Web

Ergebnisse **11 - 20** von ungefähr **2.380** für **inurl:jmx-console**. (0,15 Sekunden)

[JBoss JMX Management Console](#) - [[Diese Seite übersetzen](#)]

jboss. database=localDB,service=Hypersonic;
name=PropertyEditorManager,type=Service;
name=SystemProperties,type=Service; service=AttributePersistenceService ...

[www.crossway.com.br/jmx-console/](#) - 67k - [Zusätzliches Ergebnis](#) -

[im Cache](#) - [Ähnliche Seiten](#)

Anzeigen

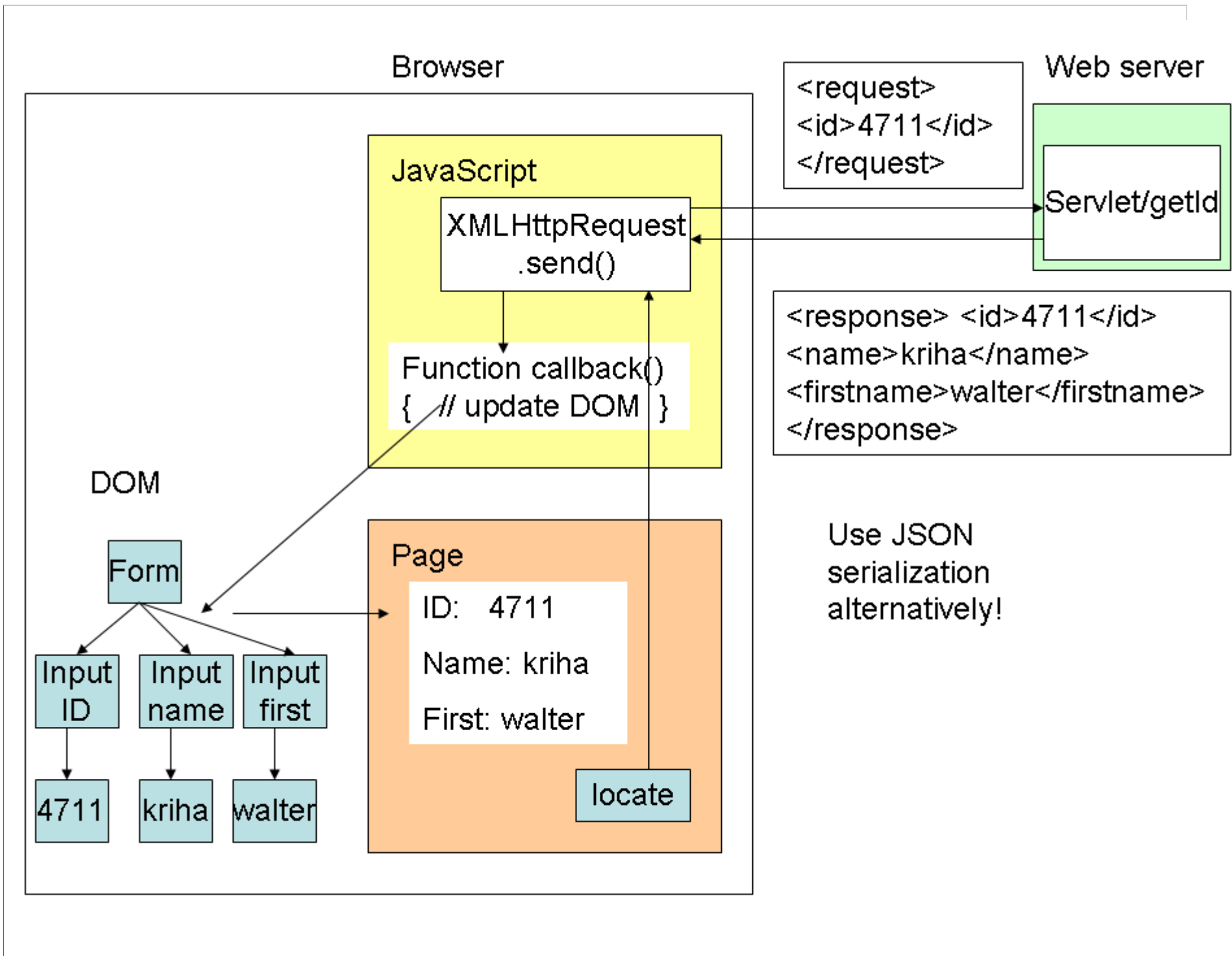
[Free JMX Consoles](#)

WebLogic JMX, MX4J, JMX1.2, JBossMX
WebSphere JMX, DBs, Systems
[manageengine.adventnet.com](#)

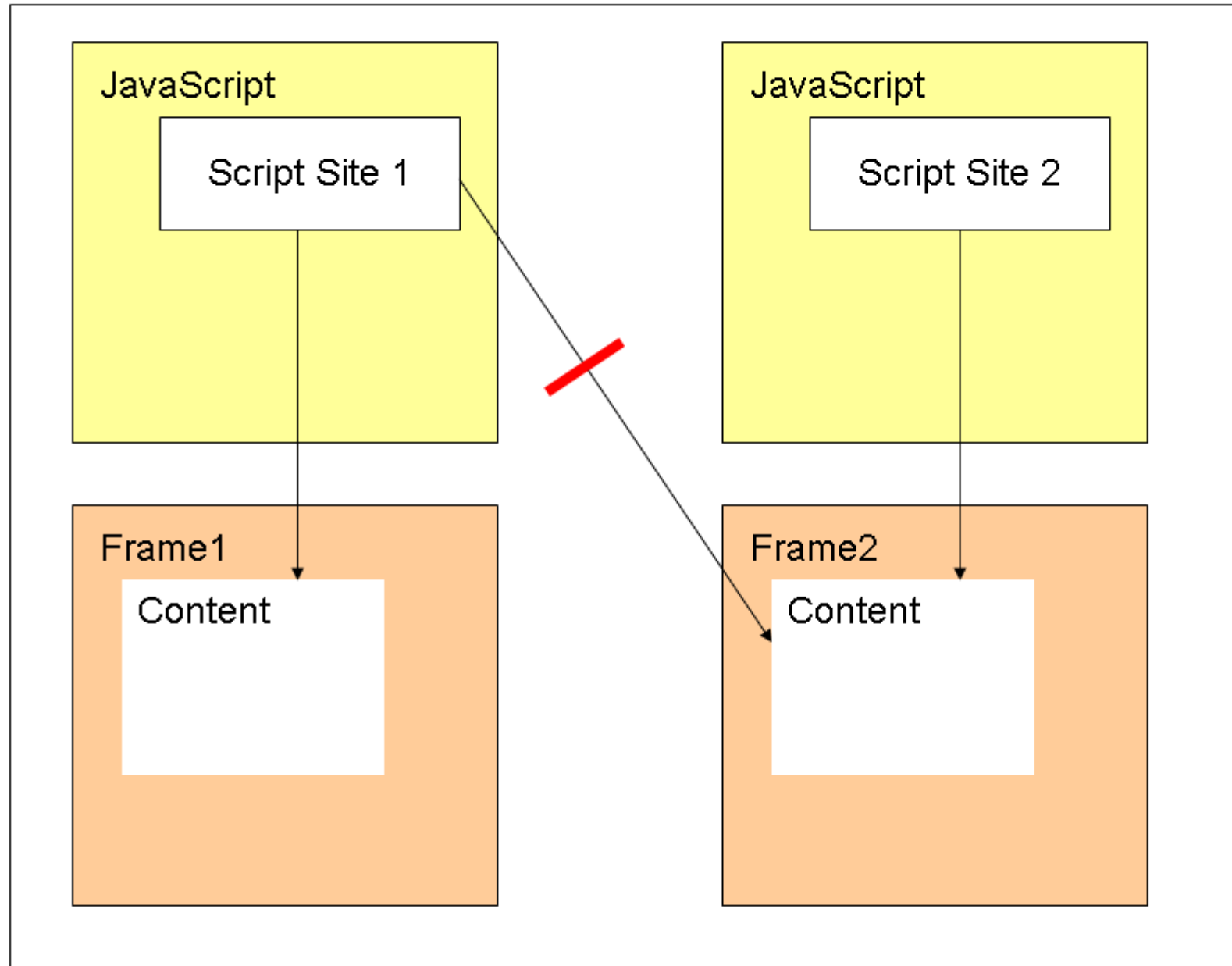
AMAZON E-Commerce Service

Sample REST Request with Style Parameter

```
http://webservices.amazon.com/onca/xml?Service=AWSECommerceService &  
AWSAccessKeyId=[Your Access Key ID Here] &  
Operation=ItemLookup &IdType=ASIN &  
ItemId=B00008OE6I &ResponseGroup=Large &  
Style=http://www.yourdomain.com/your-xsl-style-sheet.xsl
```



Page



**Web 2.0 Community
Wiki/Place Web Server**

Browser User 1

Page

ID: 4711
Name: kriha
First: walter

locate

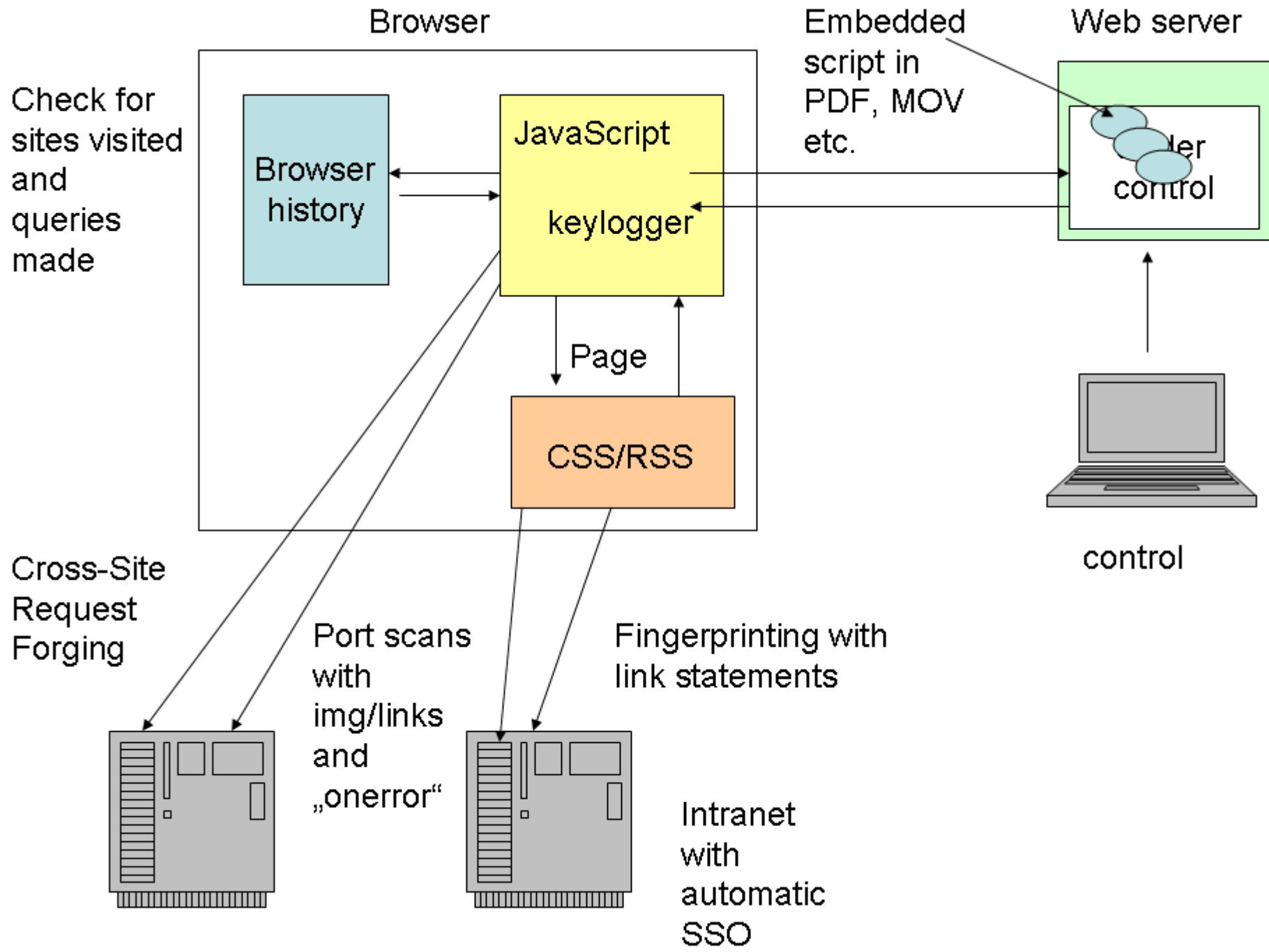
Profile User 1

Script Profile User 2

Common Pages

Common Pages

Same domain and public!



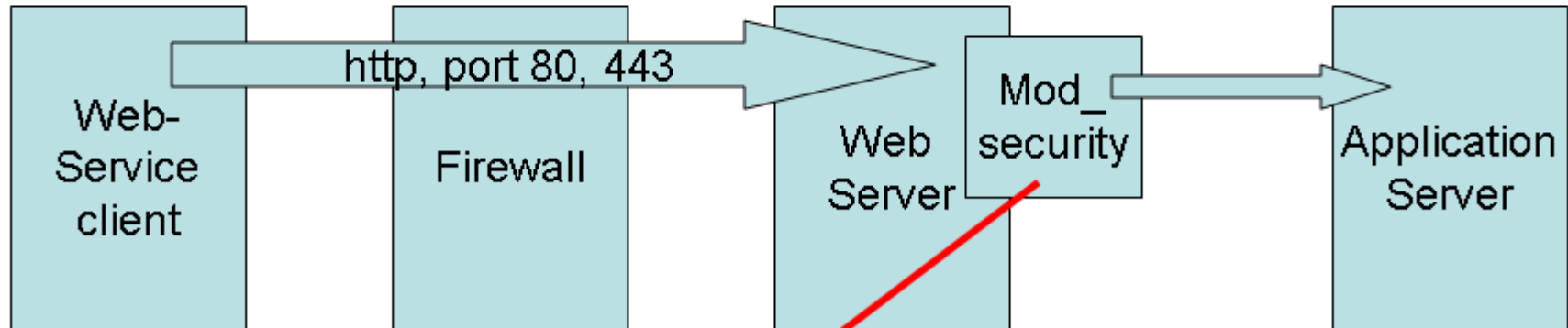
Known Mitigation Examples

WAF Filtering, Network level
filtering

SecFilterSelective Number "!^([0-9]{1,9})\$"

- Check Number for:
- Length
 - Characters/Meta
 - SQL commands

Check request for
Soap faultcode (avoid
exposure of error
information)



POST /InStock HTTP/1.1 Host: www.example.org Content-Type: application/soap+xml;
charset=utf-8 Content-Length: nnn

<?xml version="1.0"?>

<soap:Envelope xmlns:soap="http://www.w3.org/2001/12/soap-envelope"
soap:encodingStyle="http://www.w3.org/2001/12/soap-encoding">

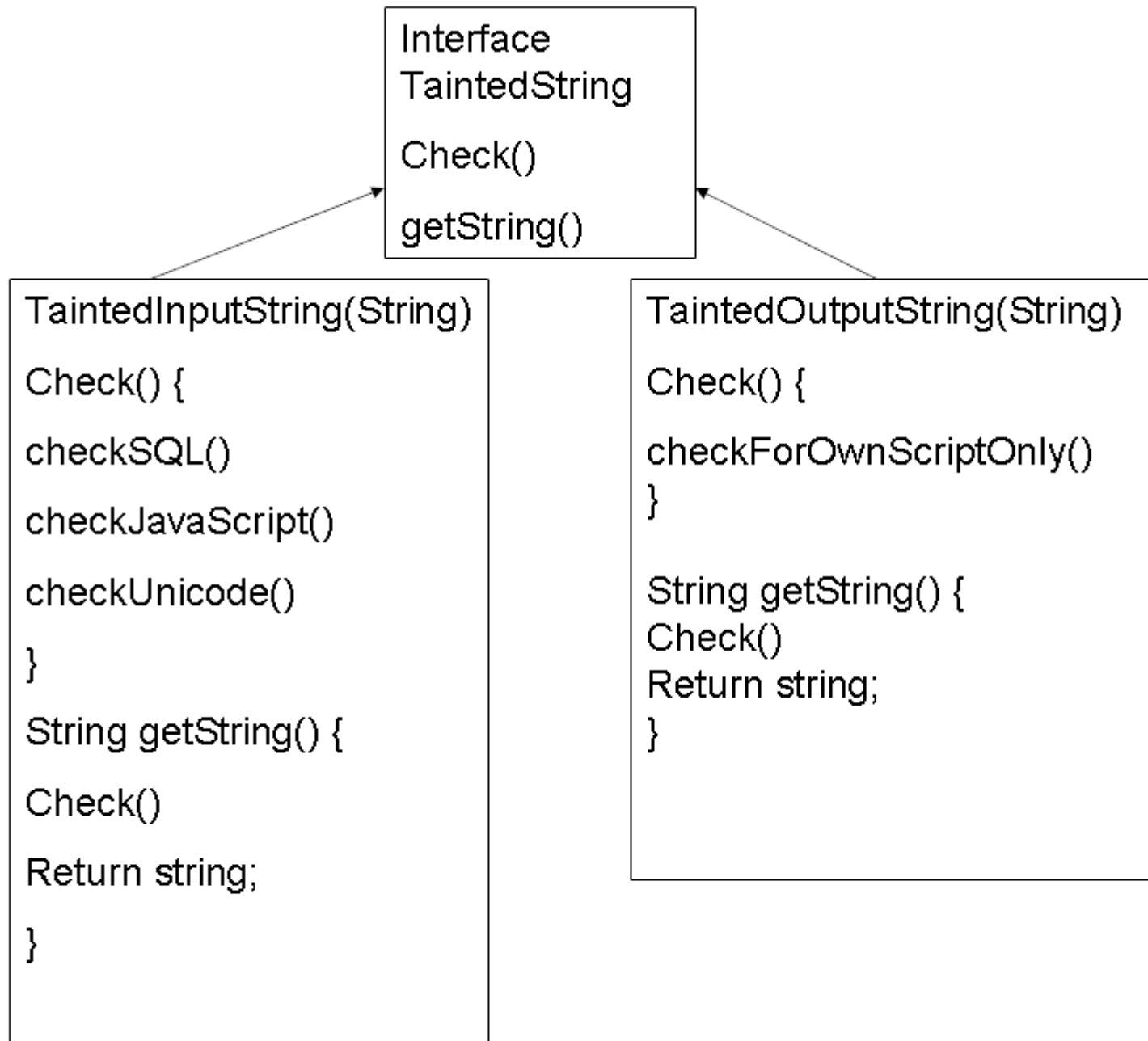
<soap:Body xmlns:k="http://www.kriha.org/number"> <m:GetId>
<m:Number>4711</m:Number> </m:GetId>

</soap:Body>

</soap:Envelope>

Other security related features of mod_security:

- URL checking
- Unicode normalization
- Message canonicalization for filtering
- Stateful filtering of selected requests
- Stateful connection of input/output values
- Stateful link/request control (did the link come from the server?)



IP Header Parameters
(e.g. protocol tcp or udp)

TCP Header Parameters
(e.g port and direction)

Rules from Firewall-Policy:

If (port = 22) &&
(protocol = TCP) &&
(NIC1-outgoing)
Action: Accept

(not real IPTABLES syntax)

ICMP Header Parameters (e.g. packet size, types)

external network address

NIC1

Packet
Packetfilter

internal network address

NIC2

destination/source address

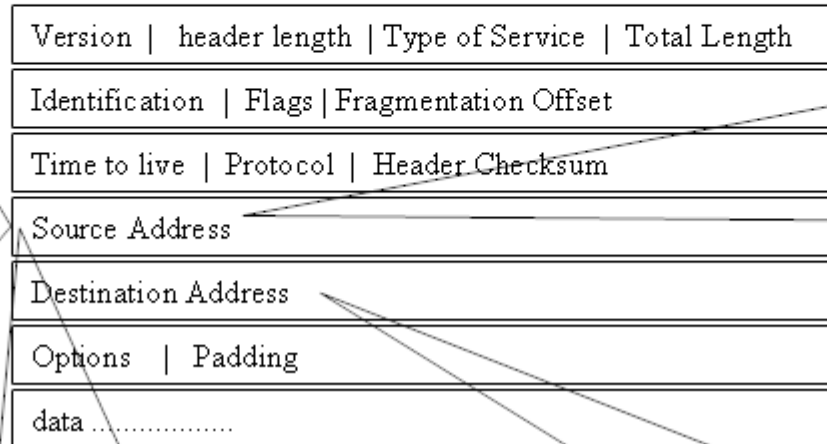
destination/source address

from : to
xxx(20) yyy(4567), tcp
yyy(4567) xxx(20), tcp

To Intranet

To Internet

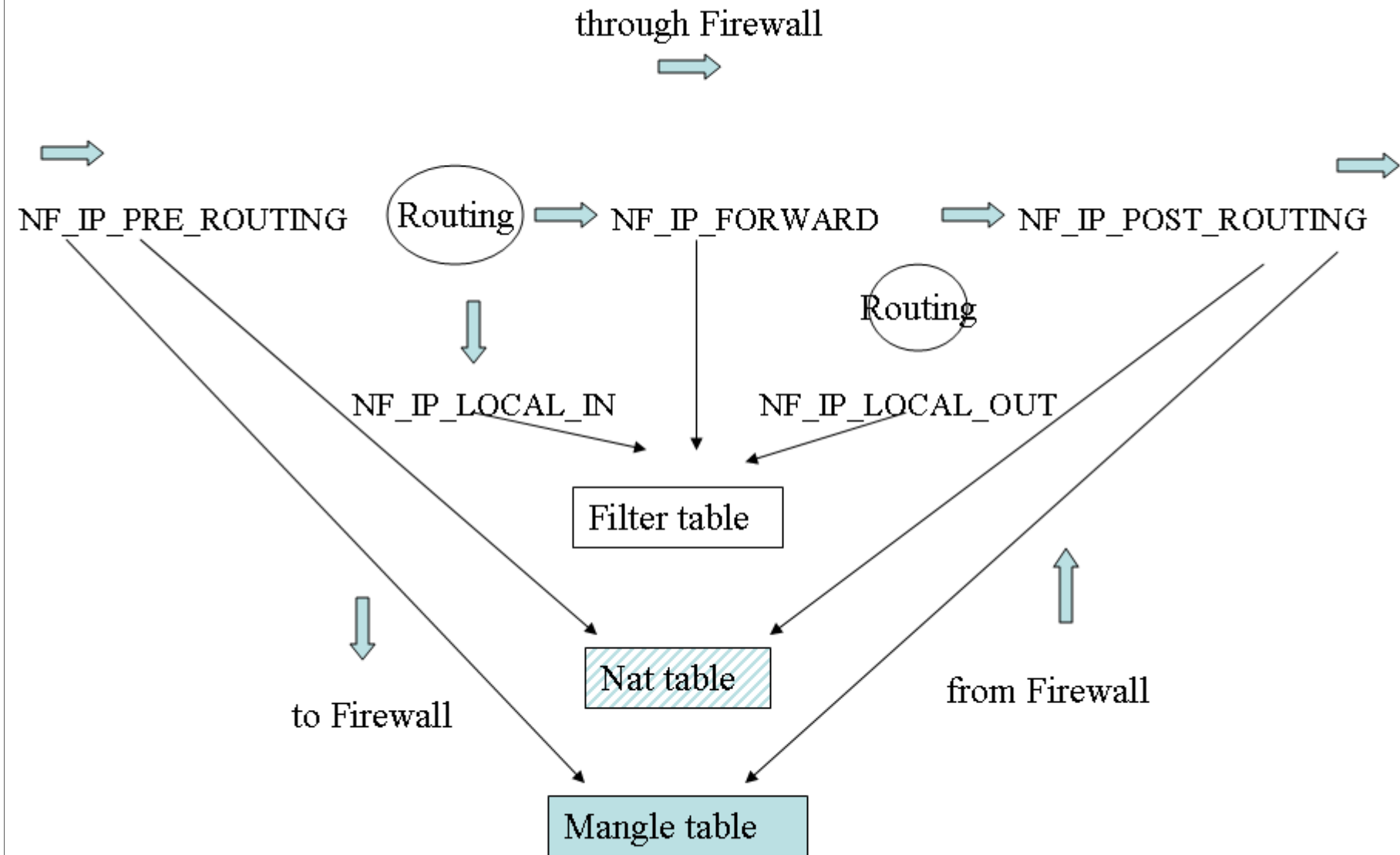
Network Address Translation (NAT) means that the source or destination address of a packet is changed



masquerading is almost like SNAT only that there is no static IP address. Instead, the source address is dynamically grabbed from an ISP, e.g via DHCP, pppoe etc.

With Source NAT (SNAT), the source address is changed, e.g. to map from private IP addresses to the real IP address of a firewall, thereby hiding the internal network.

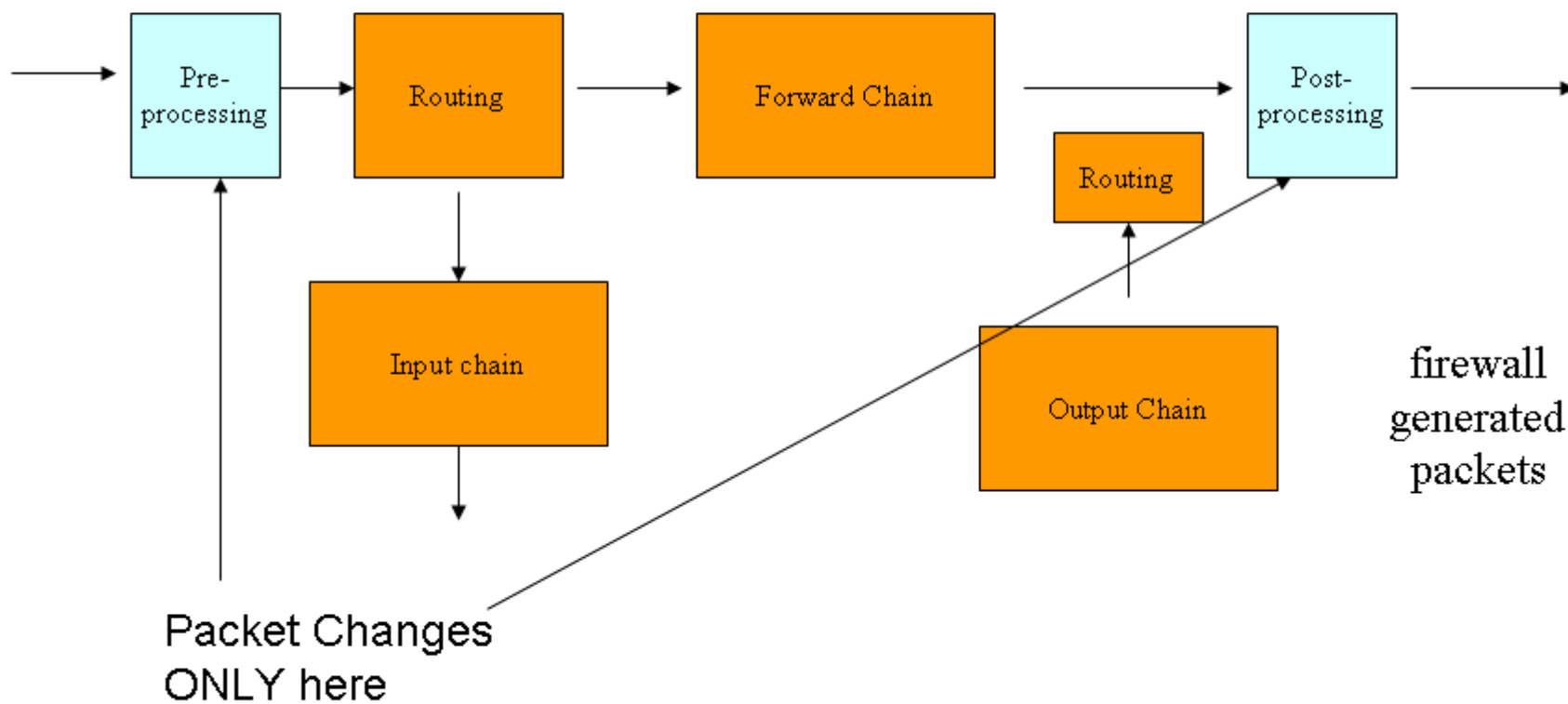
With Destination NAT (DNAT) the target address is changed, e.g. to allow transparent proxying or load-balancing



Destination NAT

all input not directed at the firewall itself goes here

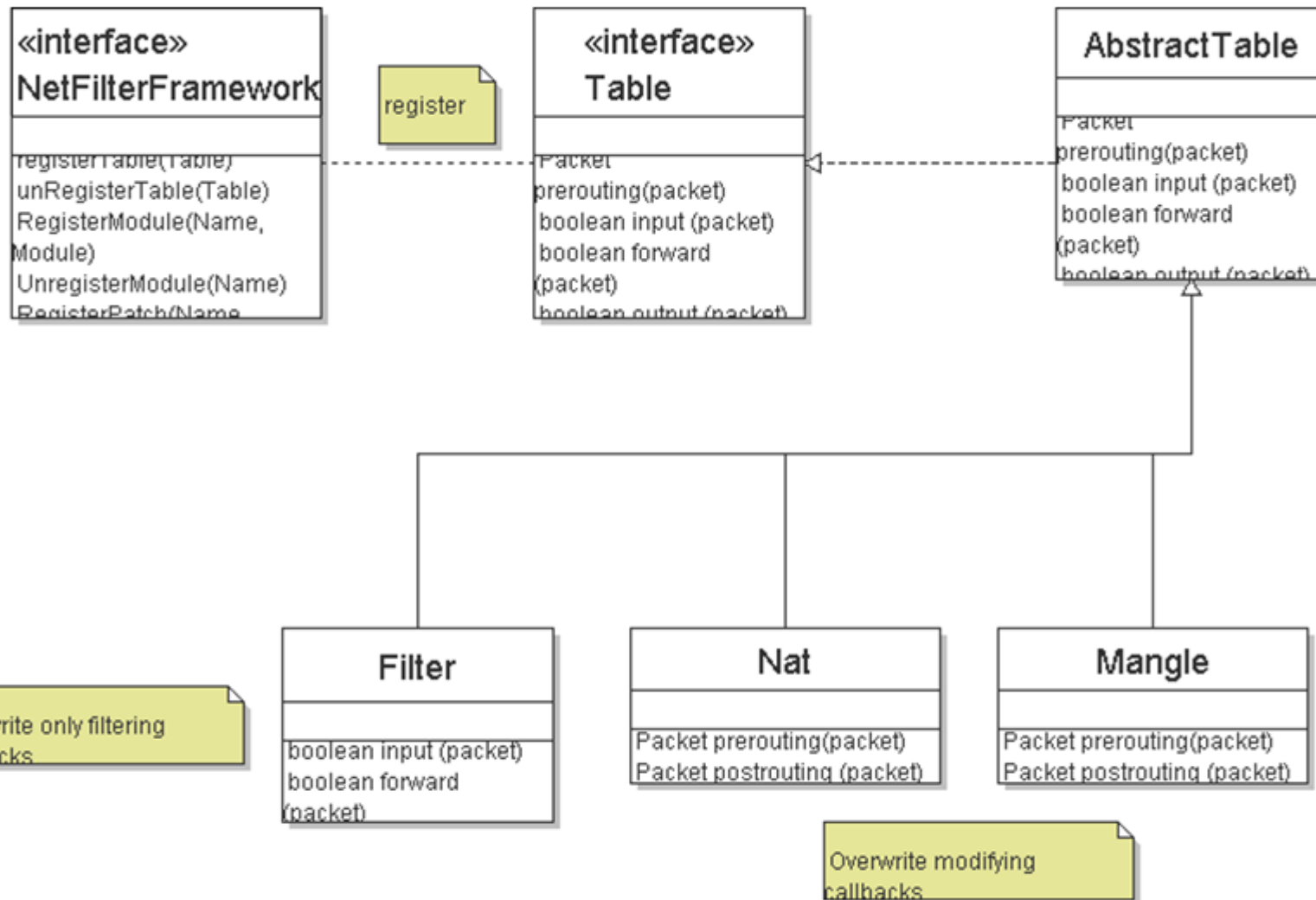
Source NAT happens here

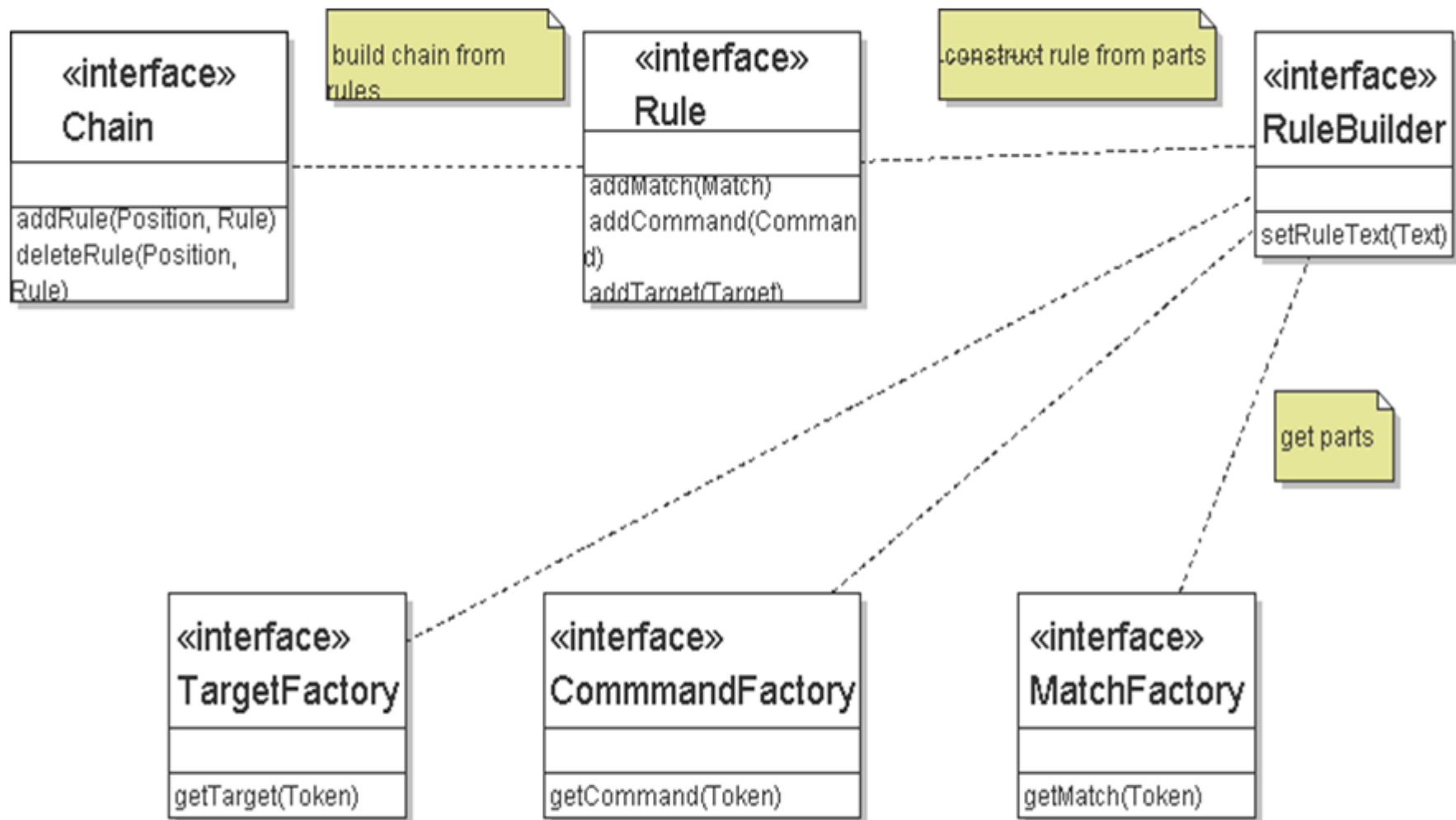


iptables -t table -command [chain] [match] -j [target/jump]

Example:

- `iptables -T FILTER -A INPUT -i $IFACE -p tcp --sport 80 -m state --state ESTABLISHED -j ACCEPT` (allow incoming web traffic if it belongs to a previous outgoing request)
- `iptables -A INPUT -i $IFACE -p tcp --sport 20 -m state --state ESTABLISHED, RELATED -j ACCEPT` (allow incoming ACTIVE ftp traffic if it belongs to a previous outgoing request, even though the incoming request is for a new – but related - port)
- `iptables -A INPUT -i $IFACE -p udp -j LOG --log-prefix „UDP Incoming:“`
- `iptables -A INPUT -i $IFACE -p udp -j DROP` (log and drop all udp traffic)



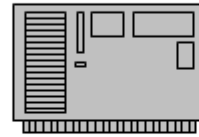


192.168.1.0/24
(intranet)

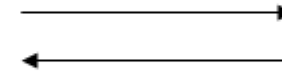
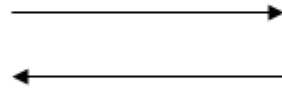
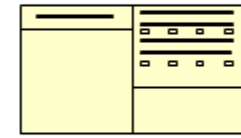


192.168.1.250

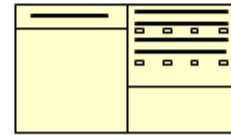
filter (firewall)



(internet)



192.84.219.128



smtp host

192.84.219.129

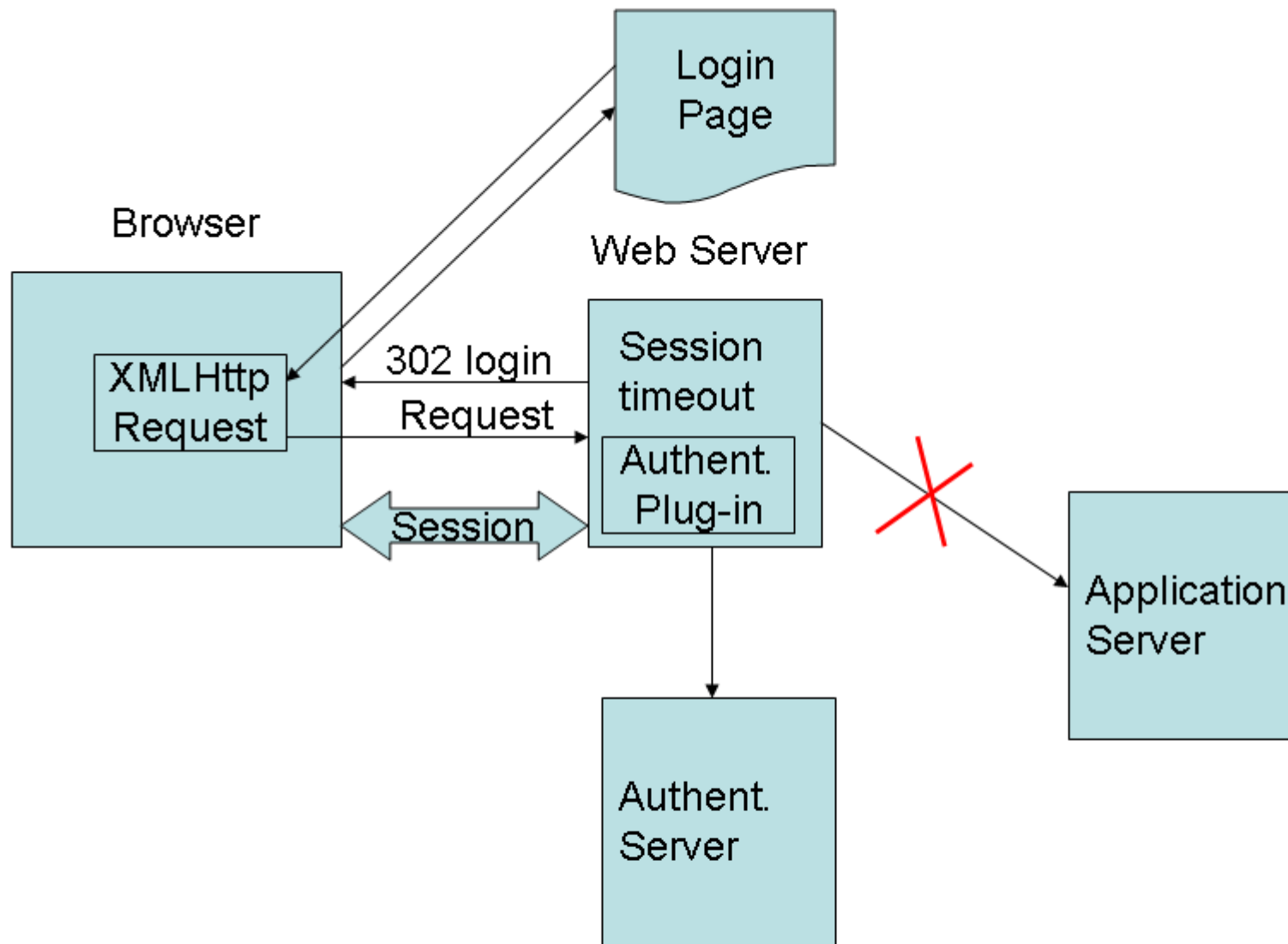


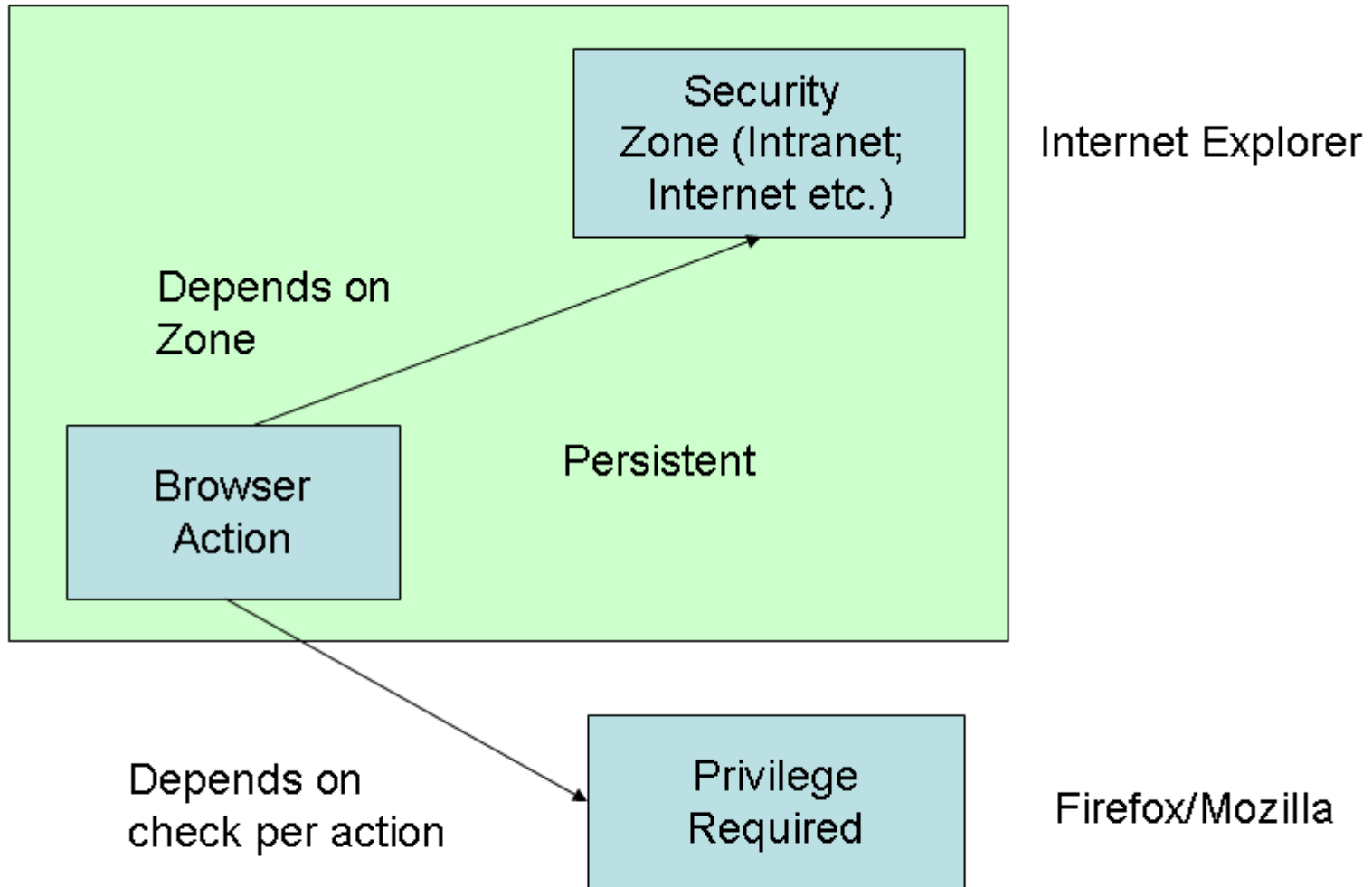
DNS host

192.84.219.130



WEB host





Fundamental Questions

Input Validation

- Are Regexp checks enough?
- How do Servlet Filters work?
- How to separate Non-terminals from terminals?
- Forwarding of modified request data – the problem of double-decoding
- Is application input a language? Of what type? How expressed? Design question?
- Tainting as a software mechanism

Filtering

- Anti-patterns of filter use?
- Proof of correctness – is illegal input blocked?
- Proof of liveness – does legal input still get through?
- Mixing of reject and accept statements?
- Filter models and automated checkers?
- Filter positions in software?

Concurrency

- Libraries for safe shell programming?
- Is shared state multithreading reliable and predictable?
- Architectures for safe concurrency (Miller)?
- Active Objects, CSP etc.

Ambient Authority

- How to restrict system call access?
- How to prevent arbitrary initial authority?
- Software architectures to achieve loader isolation?
- Language features for secure software?
- Damage control features in operating systems, languages and applications

Signs and Minds

- How to avoid confusion about identity?
- How to represent system messages reliably and without chance for fake messages?
- Software technology to establish a trusted path for users?
- Character sets and representations as fonts?
- Reliable detection of character aliases?